

# Maximal trees on $\mathcal{P}(\omega)/fin$ .

Maximal trees

Jonathan  
Cancino

## Maximal trees on $\mathcal{P}(\omega)/fin$

JONATHAN CANCINO-MANRÍQUEZ  
UNAM-UMSNH, Morelia  
México.  
jcancino@matmor.unam.mx

Winter School in Abstract Analysis, section Set Theory  
Febrero 2, 2016

This is joint work with Michael Hrušák, Gabriela Campero and Favio Miranda

## Maximal trees

Jonathan  
Cancino

# Conventions.

Maximal trees

Jonathan  
Cancino

- As usual, instead of working with  $\mathcal{P}(\omega)/fin$ , we will be working with  $[\omega]^\omega$ .
- Our trees grow downward.
- Making abuse of notation, we will make reference to trees on  $[\omega]^\omega$  as trees on  $\mathcal{P}(\omega)/fin$ , and viceverse.
- We consider  $[\omega]^\omega$  ordered by the almost containment  $\subseteq^*$ : given  $A, B \in [\omega]^\omega$ , we say that  $A \subseteq^* B$  if and only if  $A \setminus B$  is finite.

# Conventions.

Maximal trees

Jonathan  
Cancino

- As usual, instead of working with  $\mathcal{P}(\omega)/fin$ , we will be working with  $[\omega]^\omega$ .
- Our trees grow downward.
- Making abuse of notation, we will make reference to trees on  $[\omega]^\omega$  as trees on  $\mathcal{P}(\omega)/fin$ , and viceverse.
- We consider  $[\omega]^\omega$  ordered by the almost containment  $\subseteq^*$ : given  $A, B \in [\omega]^\omega$ , we say that  $A \subseteq^* B$  if and only if  $A \setminus B$  is finite.

# Conventions.

Maximal trees

Jonathan  
Cancino

- As usual, instead of working with  $\mathcal{P}(\omega)/fin$ , we will be working with  $[\omega]^\omega$ .
- Our trees grow downward.
- Making abuse of notation, we will make reference to trees on  $[\omega]^\omega$  as trees on  $\mathcal{P}(\omega)/fin$ , and viceverse.
- We consider  $[\omega]^\omega$  ordered by the almost containment  $\subseteq^*$ : given  $A, B \in [\omega]^\omega$ , we say that  $A \subseteq^* B$  if and only if  $A \setminus B$  is finite.

# Conventions.

Maximal trees

Jonathan  
Cancino

- As usual, instead of working with  $\mathcal{P}(\omega)/fin$ , we will be working with  $[\omega]^\omega$ .
- Our trees grow downward.
- Making abuse of notation, we will make reference to trees on  $[\omega]^\omega$  as trees on  $\mathcal{P}(\omega)/fin$ , and viceverse.
- We consider  $[\omega]^\omega$  ordered by the almost containment  $\subseteq^*$ : given  $A, B \in [\omega]^\omega$ , we say that  $A \subseteq^* B$  if and only if  $A \setminus B$  is finite.

# Trees on $\mathcal{P}(\omega)/fin$

Maximal trees

Jonathan  
Cancino

## Definition

A tree  $\mathcal{T}$  on  $\mathcal{P}(\omega)/fin$  is a family of elements of  $\mathcal{P}(\omega)/fin$ , such that for all  $A \in \mathcal{T}$ , the set  $pred_{\mathcal{T}}(A)$  is well ordered by  $\supseteq^*$ , the reverse ordering of  $\subseteq^*$ .

$pred_{\mathcal{T}} = \{B \in \mathcal{T} : A \subseteq^* B\}$  is the set of predecessors of  $A$  in the tree  $\mathcal{T}$ .

## Definition

Given two trees  $T, S$  on  $\mathcal{P}(\omega)/fin$ , let us say that  $T \sqsubseteq S$  if and only if  $S$  is an end extension of  $T$ , that is, for every  $x \in T$ , the sets  $pred_T(x) = pred_S(x)$ .



# Trees on $\mathcal{P}(\omega)/fin$

Maximal trees

Jonathan  
Cancino

## Definition

A tree  $\mathcal{T}$  on  $\mathcal{P}(\omega)/fin$  is a family of elements of  $\mathcal{P}(\omega)/fin$ , such that for all  $A \in \mathcal{T}$ , the set  $pred_{\mathcal{T}}(A)$  is well ordered by  $\supseteq^*$ , the reverse ordering of  $\subseteq^*$ .

$pred_{\mathcal{T}} = \{B \in \mathcal{T} : A \subseteq^* B\}$  is the set of predecessors of  $A$  in the tree  $\mathcal{T}$ .

## Definition

Given two trees  $T, S$  on  $\mathcal{P}(\omega)/fin$ , let us say that  $T \sqsubseteq S$  if and only if  $S$  is an end extension of  $T$ , that is, for every  $x \in T$ , the sets  $pred_T(x) = pred_S(x)$ .

# Trees on $\mathcal{P}(\omega)/fin$

Maximal trees

Jonathan  
Cancino

## Definition

A tree  $\mathcal{T}$  on  $\mathcal{P}(\omega)/fin$  is a family of elements of  $\mathcal{P}(\omega)/fin$ , such that for all  $A \in \mathcal{T}$ , the set  $pred_{\mathcal{T}}(A)$  is well ordered by  $\supseteq^*$ , the reverse ordering of  $\subseteq^*$ .

$pred_{\mathcal{T}} = \{B \in \mathcal{T} : A \subseteq^* B\}$  is the set of predecessors of  $A$  in the tree  $\mathcal{T}$ .

## Definition

Given two trees  $T, S$  on  $\mathcal{P}(\omega)/fin$ , let us say that  $T \sqsubseteq S$  if and only if  $S$  is an end extension of  $T$ , that is, for every  $x \in T$ , the sets  $pred_T(x) = pred_S(x)$ .

Then the set of all trees on  $\mathcal{P}(\omega)/fin$ , ordered by  $\sqsubseteq$ , satisfies the conditions of Zorn's Lemma, so this ordering has maximal elements.

#### Definition(D. Monk)

Define the cardinal invariant  $tt$  as the minimum possible size of a maximal tree on  $\mathcal{P}(\omega)/fin$ , that is,

$$tt = \min\{|\mathcal{T}| : \mathcal{T} \subseteq \mathcal{P}(\omega)/fin \text{ is a maximal tree}\}$$

Monk's notation differs from ours. Given a boolean algebra  $\mathbb{B}$  he writes  $Inc_{mm}^{tree}(\mathbb{B})$  to denote the minimum cardinality of a tree on the boolean algebra  $\mathbb{B}$ .

Then the set of all trees on  $\mathcal{P}(\omega)/fin$ , ordered by  $\sqsubseteq$ , satisfies the conditions of Zorn's Lemma, so this ordering has maximal elements.

### Definition(D. Monk)

Define the cardinal invariant  $\mathfrak{tt}$  as the minimum possible size of a maximal tree on  $\mathcal{P}(\omega)/fin$ , that is,

$$\mathfrak{tt} = \min\{|\mathcal{T}| : \mathcal{T} \subseteq \mathcal{P}(\omega)/fin \text{ is a maximal tree}\}$$

Monk's notation differs from ours. Given a boolean algebra  $\mathbb{B}$  he writes  $Inc_{mm}^{tree}(\mathbb{B})$  to denote the minimum cardinality of a tree on the boolean algebra  $\mathbb{B}$ .

Then the set of all trees on  $\mathcal{P}(\omega)/fin$ , ordered by  $\sqsubseteq$ , satisfies the conditions of Zorn's Lemma, so this ordering has maximal elements.

### Definition(D. Monk)

Define the cardinal invariant  $\mathfrak{tt}$  as the minimum possible size of a maximal tree on  $\mathcal{P}(\omega)/fin$ , that is,

$$\mathfrak{tt} = \min\{|\mathcal{T}| : \mathcal{T} \subseteq \mathcal{P}(\omega)/fin \text{ is a maximal tree}\}$$

Monk's notation differs from ours. Given a boolean algebra  $\mathbb{B}$  he writes  $Inc_{mm}^{tree}(\mathbb{B})$  to denote the minimum cardinality of a tree on the boolean algebra  $\mathbb{B}$ .

How does a maximal tree on  $\mathcal{P}(\omega)/fin$  look like?

### Lemma

A tree  $\mathcal{T} \subseteq [\omega]^\omega$  is a maximal tree if and only if for every set  $A \in [\omega]^\omega$ , one of the following holds:

- There is  $B \in \mathcal{T}$  such that  $B \subseteq^* A$ .
- There are  $B, C \in \mathcal{T}$  incomparable such that  $A \subseteq^* B \cap C$ .

How does a maximal tree on  $\mathcal{P}(\omega)/fin$  look like?

### Lemma

A tree  $\mathcal{T} \subseteq [\omega]^\omega$  is a maximal tree if and only if for every set  $A \in [\omega]^\omega$ , one of the following holds:

- There is  $B \in \mathcal{T}$  such that  $B \subseteq^* A$ .
- There are  $B, C \in \mathcal{T}$  incomparable such that  $A \subseteq^* B \cap C$ .

How does a maximal tree on  $\mathcal{P}(\omega)/fin$  look like?

### Lemma

A tree  $\mathcal{T} \subseteq [\omega]^\omega$  is a maximal tree if and only if for every set  $A \in [\omega]^\omega$ , one of the following holds:

- There is  $B \in \mathcal{T}$  such that  $B \subseteq^* A$ .
- There are  $B, C \in \mathcal{T}$  incomparable such that  $A \subseteq^* B \cap C$ .



How does a maximal tree on  $\mathcal{P}(\omega)/fin$  look like?

### Lemma

A tree  $\mathcal{T} \subseteq [\omega]^\omega$  is a maximal tree if and only if for every set  $A \in [\omega]^\omega$ , one of the following holds:

- There is  $B \in \mathcal{T}$  such that  $B \subseteq^* A$ .
- There are  $B, C \in \mathcal{T}$  incomparable such that  $A \subseteq^* B \cap C$ .

**Remark.**

If  $\mathcal{T}$  is a maximal tree on  $\mathcal{P}(\omega)/fin$ , then the following family is a reaping family

$$\mathcal{T} \cup \{\omega \setminus A : A \in \mathcal{T}\}$$

So in particular the reaping number is a lower bound for  $\mathfrak{tt}$ .

Question, D. Monk

Is  $\mathfrak{tt} = \mathfrak{c}$ ?

**Remark.**

If  $\mathcal{T}$  is a maximal tree on  $\mathcal{P}(\omega)/fin$ , then the following family is a reaping family

$$\mathcal{T} \cup \{\omega \setminus A : A \in \mathcal{T}\}$$

So in particular the reaping number is a lower bound for  $\mathfrak{tt}$ .

**Question, D. Monk**

Is  $\mathfrak{tt} = \mathfrak{c}$ ?

## Parametrized Diamond Principles

- This are guessing principles which are weakenings of the well known Jensen's diamond principle.
- For each Borel cardinal invariant corresponds a parametrized diamond principle.
- They are compatible with the negation of CH.
- They hold in many of the well known models of set theory.

The  $\diamond$ -like principle we are using is the corresponding to the cardinal invariant  $\tau_\sigma; \mathfrak{d}$ , the sequential composition of  $\tau_\sigma$  followed by  $\mathfrak{d}$ .

## Parametrized Diamond Principles

- This are guessing principles which are weakenings of the well known Jensen's diamond principle.
- For each Borel cardinal invariant corresponds a parametrized diamond principle.
  - They are compatible with the negation of CH.
  - They hold in many of the well known models of set theory.

The  $\diamond$ -like principle we are using is the corresponding to the cardinal invariant  $\tau_{\sigma; \mathfrak{d}}$ , the sequential composition of  $\tau_{\sigma}$  followed by  $\mathfrak{d}$ .

## Parametrized Diamond Principles

- This are guessing principles which are weakenings of the well known Jensen's diamond principle.
- For each Borel cardinal invariant corresponds a parametrized diamond principle.
- They are compatible with the negation of CH.
  - They hold in many of the well known models of set theory.

The  $\diamond$ -like principle we are using is the corresponding to the cardinal invariant  $\tau_\sigma; \mathfrak{d}$ , the sequential composition of  $\tau_\sigma$  followed by  $\mathfrak{d}$ .

## Parametrized Diamond Principles

- These are guessing principles which are weakenings of the well known Jensen's diamond principle.
- For each Borel cardinal invariant corresponds a parametrized diamond principle.
- They are compatible with the negation of CH.
- They hold in many of the well known models of set theory.

The  $\diamond$ -like principle we are using is the corresponding to the cardinal invariant  $\tau_\sigma; \mathfrak{d}$ , the sequential composition of  $\tau_\sigma$  followed by  $\mathfrak{d}$ .

## Parametrized Diamond Principles

- These are guessing principles which are weakenings of the well known Jensen's diamond principle.
- For each Borel cardinal invariant corresponds a parametrized diamond principle.
- They are compatible with the negation of CH.
- They hold in many of the well known models of set theory.

The  $\diamond$ -like principle we are using is the corresponding to the cardinal invariant  $\tau_\sigma; \mathfrak{d}$ , the sequential composition of  $\tau_\sigma$  followed by  $\mathfrak{d}$ .



## Parametrized Diamond Principles

- These are guessing principles which are weakenings of the well known Jensen's diamond principle.
- For each Borel cardinal invariant corresponds a parametrized diamond principle.
- They are compatible with the negation of CH.
- They hold in many of the well known models of set theory.

The  $\diamond$ -like principle we are using is the corresponding to the cardinal invariant  $\tau_\sigma; \mathfrak{d}$ , the sequential composition of  $\tau_\sigma$  followed by  $\mathfrak{d}$ .

We have found two different shapes for these kind of trees:

### Theorem

$\diamond_{L(\mathbb{R})}(\mathfrak{r}_\sigma; \mathfrak{d})$  implies that there is a maximal tree on  $\mathcal{P}(\omega)/fin$  of cardinality  $\omega_1$ , has height  $\omega_1$ , and all nodes, except the root of the tree (who has  $\omega_1$  successors), have exactly one successor.

### Theorem

$\diamond_{L(\mathbb{R})}(\mathfrak{r}_\sigma; \mathfrak{d})$  implies that there is a maximal tree  $\mathcal{T}$  on  $\mathcal{P}(\omega)/fin$ , such that every node  $A \in \mathcal{T}$  has  $\omega_1$  successors, and the height of  $\mathcal{T}$  is  $\omega$ .

### Corollary

In the Sacks model  $\mathfrak{tt}$  is  $\omega_1$ , while the continuum is  $\omega_2$ .

We have found two different shapes for these kind of trees:

### Theorem

$\diamond_{L(\mathbb{R})}(\mathfrak{t}_\sigma; \mathfrak{d})$  implies that there is a maximal tree on  $\mathcal{P}(\omega)/fin$  of cardinality  $\omega_1$ , has height  $\omega_1$ , and all nodes, except the root of the tree (who has  $\omega_1$  successors), have exactly one successor.

### Theorem

$\diamond_{L(\mathbb{R})}(\mathfrak{t}_\sigma; \mathfrak{d})$  implies that there is a maximal tree  $\mathcal{T}$  on  $\mathcal{P}(\omega)/fin$ , such that every node  $A \in \mathcal{T}$  has  $\omega_1$  successors, and the height of  $\mathcal{T}$  is  $\omega$ .

### Corollary

In the Sacks model  $\mathfrak{t}$  is  $\omega_1$ , while the continuum is  $\omega_2$ .

We have found two different shapes for these kind of trees:

### Theorem

$\diamond_{L(\mathbb{R})}(\mathfrak{r}_\sigma; \mathfrak{d})$  implies that there is a maximal tree on  $\mathcal{P}(\omega)/fin$  of cardinality  $\omega_1$ , has height  $\omega_1$ , and all nodes, except the root of the tree (who has  $\omega_1$  successors), have exactly one successor.

### Theorem

$\diamond_{L(\mathbb{R})}(\mathfrak{r}_\sigma; \mathfrak{d})$  implies that there is a maximal tree  $\mathcal{T}$  on  $\mathcal{P}(\omega)/fin$ , such that every node  $A \in \mathcal{T}$  has  $\omega_1$  successors, and the height of  $\mathcal{T}$  is  $\omega$ .

### Corollary

In the Sacks model  $\mathfrak{tt}$  is  $\omega_1$ , while the continuum is  $\omega_2$ .

We have found two different shapes for these kind of trees:

### Theorem

$\diamond_{L(\mathbb{R})}(\mathfrak{r}_\sigma; \mathfrak{d})$  implies that there is a maximal tree on  $\mathcal{P}(\omega)/fin$  of cardinality  $\omega_1$ , has height  $\omega_1$ , and all nodes, except the root of the tree (who has  $\omega_1$  successors), have exactly one successor.

### Theorem

$\diamond_{L(\mathbb{R})}(\mathfrak{r}_\sigma; \mathfrak{d})$  implies that there is a maximal tree  $\mathcal{T}$  on  $\mathcal{P}(\omega)/fin$ , such that every node  $A \in \mathcal{T}$  has  $\omega_1$  successors, and the height of  $\mathcal{T}$  is  $\omega$ .

### Corollary

In the Sacks model  $\mathfrak{tt}$  is  $\omega_1$ , while the continuum is  $\omega_2$ .

In the construction of the trees in the two theorems, we make use of the dominating number, and it is not clear how to skip this, so one may ask whether the dominating number  $\mathfrak{d}$  is a lower bound of  $\mathfrak{t}$ .

### Question

Is  $\mathfrak{d}$  a lower bound for  $\mathfrak{t}$ ?

We only have partial evidence about this.

In the construction of the trees in the two theorems, we make use of the dominating number, and it is not clear how to skip this, so one may ask whether the dominating number  $\mathfrak{d}$  is a lower bound of  $\mathfrak{tt}$ .

### Question

Is  $\mathfrak{d}$  a lower bound for  $\mathfrak{tt}$ ?

We only have partial evidence about this.

In the construction of the trees in the two theorems, we make use of the dominating number, and it is not clear how to skip this, so one may ask whether the dominating number  $\mathfrak{d}$  is a lower bound of  $\mathfrak{t}$ .

### Question

Is  $\mathfrak{d}$  a lower bound for  $\mathfrak{t}$ ?

We only have partial evidence about this.



## Definition

Let  $\mathcal{T}$  be a tree on  $\mathcal{P}(\omega)/fin$ . We say that the tree  $\mathcal{T}$  is an ideal-tree if for every  $A \in \mathcal{T}$ , the family of sets  $\{A \cap B : B \notin pred_{\mathcal{T}}(A)\}$  generates a proper ideal on  $A$ .

The tree of height  $\omega$  mentioned above is actually an ideal-tree.

## Proposition

Let  $\mathcal{T}$  be a maximal tree on  $\mathcal{P}(\omega)/fin$ . Then

- If  $\mathcal{T}$  is an ideal-tree, then it has size at least  $\mathfrak{d}$ .
- If  $\mathcal{T}$  has a branch of countable cofinality, then  $\mathfrak{d} \leq |\mathcal{T}|$ .
- If  $\mathcal{T}$  has an infinite AD family then  $\mathfrak{d} \leq |\mathcal{T}|$ .
- If  $\mathcal{T}$  has a terminal node (a node with no successors), then  $|\mathcal{T}| = \mathfrak{c}$ .

## Definition

Let  $\mathcal{T}$  be a tree on  $\mathcal{P}(\omega)/fin$ . We say that the tree  $\mathcal{T}$  is an ideal-tree if for every  $A \in \mathcal{T}$ , the family of sets  $\{A \cap B : B \notin pred_{\mathcal{T}}(A)\}$  generates a proper ideal on  $A$ .

The tree of height  $\omega$  mentioned above is actually an ideal-tree.

## Proposition

Let  $\mathcal{T}$  be a maximal tree on  $\mathcal{P}(\omega)/fin$ . Then

- If  $\mathcal{T}$  is an ideal-tree, then it has size at least  $\mathfrak{d}$ .
- If  $\mathcal{T}$  has a branch of countable cofinality, then  $\mathfrak{d} \leq |\mathcal{T}|$ .
- If  $\mathcal{T}$  has an infinite AD family then  $\mathfrak{d} \leq |\mathcal{T}|$ .
- If  $\mathcal{T}$  has a terminal node (a node with no successors), then  $|\mathcal{T}| = \mathfrak{c}$ .

## Definition

Let  $\mathcal{T}$  be a tree on  $\mathcal{P}(\omega)/fin$ . We say that the tree  $\mathcal{T}$  is an ideal-tree if for every  $A \in \mathcal{T}$ , the family of sets  $\{A \cap B : B \notin pred_{\mathcal{T}}(A)\}$  generates a proper ideal on  $A$ .

The tree of height  $\omega$  mentioned above is actually an ideal-tree.

## Proposition

Let  $\mathcal{T}$  be a maximal tree on  $\mathcal{P}(\omega)/fin$ . Then

- If  $\mathcal{T}$  is an ideal-tree, then it has size at least  $\mathfrak{d}$ .
- If  $\mathcal{T}$  has a branch of countable cofinality, then  $\mathfrak{d} \leq |\mathcal{T}|$ .
- If  $\mathcal{T}$  has an infinite AD family then  $\mathfrak{d} \leq |\mathcal{T}|$ .
- If  $\mathcal{T}$  has a terminal node (a node with no successors), then  $|\mathcal{T}| = \mathfrak{c}$ .

## Definition

Let  $\mathcal{T}$  be a tree on  $\mathcal{P}(\omega)/fin$ . We say that the tree  $\mathcal{T}$  is an ideal-tree if for every  $A \in \mathcal{T}$ , the family of sets  $\{A \cap B : B \notin pred_{\mathcal{T}}(A)\}$  generates a proper ideal on  $A$ .

The tree of height  $\omega$  mentioned above is actually an ideal-tree.

## Proposition

Let  $\mathcal{T}$  be a maximal tree on  $\mathcal{P}(\omega)/fin$ . Then

- If  $\mathcal{T}$  is an ideal-tree, then it has size at least  $\mathfrak{d}$ .
- If  $\mathcal{T}$  has a branch of countable cofinality, then  $\mathfrak{d} \leq |\mathcal{T}|$ .
- If  $\mathcal{T}$  has an infinite AD family then  $\mathfrak{d} \leq |\mathcal{T}|$ .
- If  $\mathcal{T}$  has a terminal node (a node with no successors), then  $|\mathcal{T}| = \mathfrak{c}$ .

## Definition

Let  $\mathcal{T}$  be a tree on  $\mathcal{P}(\omega)/fin$ . We say that the tree  $\mathcal{T}$  is an ideal-tree if for every  $A \in \mathcal{T}$ , the family of sets  $\{A \cap B : B \notin pred_{\mathcal{T}}(A)\}$  generates a proper ideal on  $A$ .

The tree of height  $\omega$  mentioned above is actually an ideal-tree.

## Proposition

Let  $\mathcal{T}$  be a maximal tree on  $\mathcal{P}(\omega)/fin$ . Then

- If  $\mathcal{T}$  is an ideal-tree, then it has size at least  $\mathfrak{d}$ .
- If  $\mathcal{T}$  has a branch of countable cofinality, then  $\mathfrak{d} \leq |\mathcal{T}|$ .
- If  $\mathcal{T}$  has an infinite *AD* family then  $\mathfrak{d} \leq |\mathcal{T}|$ .
- If  $\mathcal{T}$  has a terminal node (a node with no successors), then  $|\mathcal{T}| = \mathfrak{c}$ .

## Definition

Let  $\mathcal{T}$  be a tree on  $\mathcal{P}(\omega)/fin$ . We say that the tree  $\mathcal{T}$  is an ideal-tree if for every  $A \in \mathcal{T}$ , the family of sets  $\{A \cap B : B \notin pred_{\mathcal{T}}(A)\}$  generates a proper ideal on  $A$ .

The tree of height  $\omega$  mentioned above is actually an ideal-tree.

## Proposition

Let  $\mathcal{T}$  be a maximal tree on  $\mathcal{P}(\omega)/fin$ . Then

- If  $\mathcal{T}$  is an ideal-tree, then it has size at least  $\mathfrak{d}$ .
- If  $\mathcal{T}$  has a branch of countable cofinality, then  $\mathfrak{d} \leq |\mathcal{T}|$ .
- If  $\mathcal{T}$  has an infinite *AD* family then  $\mathfrak{d} \leq |\mathcal{T}|$ .
- If  $\mathcal{T}$  has a terminal node (a node with no sucesors), then  $|\mathcal{T}| = \mathfrak{c}$ .

## Definition

Let  $\mathcal{T}$  be a tree on  $\mathcal{P}(\omega)/fin$ . We say that the tree  $\mathcal{T}$  is an ideal-tree if for every  $A \in \mathcal{T}$ , the family of sets  $\{A \cap B : B \notin pred_{\mathcal{T}}(A)\}$  generates a proper ideal on  $A$ .

The tree of height  $\omega$  mentioned above is actually an ideal-tree.

## Proposition

Let  $\mathcal{T}$  be a maximal tree on  $\mathcal{P}(\omega)/fin$ . Then

- If  $\mathcal{T}$  is an ideal-tree, then it has size at least  $\mathfrak{d}$ .
- If  $\mathcal{T}$  has a branch of countable cofinality, then  $\mathfrak{d} \leq |\mathcal{T}|$ .
- If  $\mathcal{T}$  has an infinite *AD* family then  $\mathfrak{d} \leq |\mathcal{T}|$ .
- If  $\mathcal{T}$  has a terminal node (a node with no successors), then  $|\mathcal{T}| = \mathfrak{c}$ .

## Theorem

It is consistent that  $\mathfrak{tr} < \text{non}(\mathcal{M})$ . In particular it is consistent  $\mathfrak{tr} < \mathfrak{i}$ .

Guideline of proof:

- Make a  $\omega_2$ -length CSI of any of your favourite fat tree forcing.
- This forcing is  $\omega^\omega$ -bounding and adds eventually different reals.
- It was proved by J. Zapletal that this forcing preserves Ramsey ultrafilters.

So in the final extension the  $\sigma$ -reaping number and the dominating number are both  $\omega_1$ , meanwhile  $\text{non}(\mathcal{M})$  is big. Since this forcing is a definable forcing notion, it follows that  $\diamond_{L(\mathbb{R})}(\mathfrak{r}_\sigma, \mathfrak{d})$  holds.



## Theorem

It is consistent that  $\mathfrak{tr} < \text{non}(\mathcal{M})$ . In particular it is consistent  $\mathfrak{tr} < \mathfrak{i}$ .

Guideline of proof:

- Make a  $\omega_2$ -length CSI of any of your favourite fat tree forcing.
- This forcing is  $\omega^\omega$ -bounding and adds eventually different reals.
- It was proved by J. Zapletal that this forcing preserves Ramsey ultrafilters.

So in the final extension the  $\sigma$ -reaping number and the dominating number are both  $\omega_1$ , meanwhile  $\text{non}(\mathcal{M})$  is big. Since this forcing is a definable forcing notion, it follows that  $\diamond_{L(\mathbb{R})}(\mathfrak{r}_\sigma, \mathfrak{d})$  holds.

## Theorem

It is consistent that  $\mathfrak{tr} < \text{non}(\mathcal{M})$ . In particular it is consistent  $\mathfrak{tr} < \mathfrak{i}$ .

Guideline of proof:

- Make a  $\omega_2$ -length CSI of any of your favourite fat tree forcing.
- This forcing is  $\omega^\omega$ -bounding and adds eventually different reals.
- It was proved by J. Zapletal that this forcing preserves Ramsey ultrafilters.

So in the final extension the  $\sigma$ -reaping number and the dominating number are both  $\omega_1$ , meanwhile  $\text{non}(\mathcal{M})$  is big. Since this forcing is a definable forcing notion, it follows that  $\diamond_{L(\mathbb{R})}(\mathfrak{r}_\sigma, \mathfrak{d})$  holds.

## Theorem

It is consistent that  $\mathfrak{tr} < \text{non}(\mathcal{M})$ . In particular it is consistent  $\mathfrak{tr} < \mathfrak{i}$ .

Guideline of proof:

- Make a  $\omega_2$ -length CSI of any of your favourite fat tree forcing.
- This forcing is  $\omega^\omega$ -bounding and adds eventually different reals.
- It was proved by J. Zapletal that this forcing preserves Ramsey ultrafilters.

So in the final extension the  $\sigma$ -reaping number and the dominating number are both  $\omega_1$ , meanwhile  $\text{non}(\mathcal{M})$  is big. Since this forcing is a definable forcing notion, it follows that  $\diamond_{L(\mathbb{R})}(\mathfrak{r}_\sigma, \mathfrak{d})$  holds.

## Theorem

It is consistent that  $\mathfrak{tr} < \text{non}(\mathcal{M})$ . In particular it is consistent  $\mathfrak{tr} < \mathfrak{i}$ .

Guideline of proof:

- Make a  $\omega_2$ -length CSI of any of your favourite fat tree forcing.
- This forcing is  $\omega^\omega$ -bounding and adds eventually different reals.
- It was proved by J. Zapletal that this forcing preserves Ramsey ultrafilters.

So in the final extension the  $\sigma$ -reaping number and the dominating number are both  $\omega_1$ , meanwhile  $\text{non}(\mathcal{M})$  is big. Since this forcing is a definable forcing notion, it follows that  $\diamond_{L(\mathbb{R})}(\mathfrak{r}_\sigma, \mathfrak{d})$  holds.

## Theorem

It is consistent that  $\mathfrak{tr} < \text{non}(\mathcal{M})$ . In particular it is consistent  $\mathfrak{tr} < \mathfrak{i}$ .

Guideline of proof:

- Make a  $\omega_2$ -length CSI of any of your favourite fat tree forcing.
- This forcing is  $\omega^\omega$ -bounding and adds eventually different reals.
- It was proved by J. Zapletal that this forcing preserves Ramsey ultrafilters.

So in the final extension the  $\sigma$ -reaping number and the dominating number are both  $\omega_1$ , meanwhile  $\text{non}(\mathcal{M})$  is big. Since this forcing is a definable forcing notion, it follows that  $\diamond_{L(\mathbb{R})}(\mathfrak{r}_\sigma, \mathfrak{d})$  holds.

Thank you for your attention!