

REVERSIBILITY OF RELATIONAL STRUCTURES

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2nd February 2016

Plan of the presentation

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- Preliminaries

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- Strongly reversible, reversible and weakly reversible interpretations

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- Strongly reversible, reversible and weakly reversible interpretations
- Characterization of strongly reversible interpretations

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	$\neg(\text{b1}) (\top)$	(b1)	(b2)	(b3)	(b4)
$\neg(\text{a1})$ (\top)	$\bigcup_{\omega} \mathbb{L}_2 \cup \bigcup_{\omega} \mathbb{L}_2^0 \cup \bigcup_{\omega} \mathbb{L}_4$	$\bigcup_{\omega} \mathbb{L}_2 \cup \bigcup_{\omega} \mathbb{L}_4 \cup \mathbb{L}_4^0$	$\mathbb{L}_2 \cup \bigcup_{\omega} \mathbb{L}_2^0 \cup \bigcup_{\omega} \mathbb{L}_4^{0,2}$	$\bigcup_{\omega} \mathbb{L}_2^0 \cup \mathbb{L}_2^1 \cup \bigcup_{\omega} \mathbb{L}_4^{0,2}$	$\bigcup_{\omega} \mathbb{L}_2 \cup \bigcup_{\omega} \mathbb{L}_4$
(a1)	$\mathbb{L}_2 \cup \bigcup_{\omega} \mathbb{L}_{\omega} \cup \bigcup_{\omega} \mathbb{L}_{\omega}^0$	$\bigcup_{\omega} \mathbb{L}_2^0 \cup \bigcup_{\omega} \mathbb{L}_4^0 \cup \mathbb{L}_4^{0,2}$	$\mathbb{L}_2 \cup \bigcup_{\omega} \mathbb{L}_2^0 \cup \bigcup_{\omega} \mathbb{L}_4^0$	$\bigcup_{\omega} \mathbb{L}_2^0 \cup \mathbb{L}_2^1 \cup \bigcup_{\omega} \mathbb{L}_4^{0,3}$	$\mathbb{L}_2 \cup \bigcup_{\omega} \mathbb{L}_{\omega}$
(a2)	$\bigcup_{\omega} \mathbb{L}_{\omega} \cup \bigcup_{\omega} \mathbb{L}_{\omega}^0$	$\bigcup_{\omega} \mathbb{L}_{\omega} \cup \mathbb{L}_{\omega}^0$	$\mathbb{L}_{\omega} \cup \bigcup_{\omega} \mathbb{L}_{\omega}^0$	$\bigcup_{\omega} \mathbb{L}_{\omega}^0 \cup \mathbb{L}_{\omega}^1$	$\bigcup_{\omega} \mathbb{L}_{\omega}$
(a3)	$\bigcup_{\omega} \mathbb{L}_{\omega}^1 \cup \bigcup_{\omega} \mathbb{L}_{\omega}^{0,1}$	$\bigcup_{\omega} \mathbb{L}_{\omega}^1 \cup \mathbb{L}_{\omega}^{0,1}$	$\mathbb{L}_{\omega}^1 \cup \bigcup_{\omega} \mathbb{L}_{\omega}^{0,1}$	$\mathbb{L}_{\omega}^{0,2} \cup \bigcup_{\omega} \mathbb{L}_{\omega}^{1,2}$	$\bigcup_{\omega} \mathbb{L}_{\omega}^0$
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(a5)	$\bigcup_{\omega} \mathbb{L}_2 \cup \bigcup_{\omega} \mathbb{L}_2^0 \cup \mathbb{L}_3$	$\bigcup_{\omega} \mathbb{L}_2 \cup \mathbb{L}_2^0 \cup \mathbb{L}_3$	$\mathbb{L}_2 \cup \bigcup_{\omega} \mathbb{L}_2^0 \cup \mathbb{L}_3$	$\bigcup_{\omega} \mathbb{L}_2^0 \cup \mathbb{L}_2^1 \cup \mathbb{L}_3^0$	$\bigcup_{\omega} \mathbb{L}_2 \cup \mathbb{L}_3$

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	\neg (b1) (\top)	(b1)	(b2)	(b3)	(b4)
\neg (a1) (\top)	$\bigcup_{\omega} \mathbb{L}_2 \cup \bigcup_{\omega} \mathbb{L}_2^0 \cup \bigcup_{\omega} \mathbb{L}_4$	$\bigcup_{\omega} \mathbb{L}_2 \cup \bigcup_{\omega} \mathbb{L}_4 \cup \mathbb{L}_4^0$	$\mathbb{L}_2 \cup \bigcup_{\omega} \mathbb{L}_2^0 \cup \bigcup_{\omega} \mathbb{L}_4^{0,2}$	$\bigcup_{\omega} \mathbb{L}_2^0 \cup \mathbb{L}_2^1 \cup \bigcup_{\omega} \mathbb{L}_4^{0,2}$	$\bigcup_{\omega} \mathbb{L}_2 \cup \bigcup_{\omega} \mathbb{L}_4$
(a1)	$\mathbb{L}_2 \cup \bigcup_{\omega} \mathbb{L}_{\omega} \cup \bigcup_{\omega} \mathbb{L}_{\omega}^0$	$\bigcup_{\omega} \mathbb{L}_2^0 \cup \bigcup_{\omega} \mathbb{L}_4^0 \cup \mathbb{L}_4^{0,2}$	$\mathbb{L}_2 \cup \bigcup_{\omega} \mathbb{L}_2^0 \cup \bigcup_{\omega} \mathbb{L}_4^0$	$\bigcup_{\omega} \mathbb{L}_2^0 \cup \mathbb{L}_2^1 \cup \bigcup_{\omega} \mathbb{L}_4^{0,3}$	$\mathbb{L}_2 \cup \bigcup_{\omega} \mathbb{L}_{\omega}$
(a2)	$\bigcup_{\omega} \mathbb{L}_{\omega} \cup \bigcup_{\omega} \mathbb{L}_{\omega}^0$	$\bigcup_{\omega} \mathbb{L}_{\omega} \cup \mathbb{L}_{\omega}^0$	$\mathbb{L}_{\omega} \cup \bigcup_{\omega} \mathbb{L}_{\omega}^0$	$\bigcup_{\omega} \mathbb{L}_{\omega}^0 \cup \mathbb{L}_{\omega}^1$	$\bigcup_{\omega} \mathbb{L}_{\omega}$
(a3)	$\bigcup_{\omega} \mathbb{L}_{\omega}^1 \cup \bigcup_{\omega} \mathbb{L}_{\omega}^{0,1}$	$\bigcup_{\omega} \mathbb{L}_{\omega}^1 \cup \mathbb{L}_{\omega}^{0,1}$	$\mathbb{L}_{\omega}^1 \cup \bigcup_{\omega} \mathbb{L}_{\omega}^{0,1}$	$\mathbb{L}_{\omega}^{0,2} \cup \bigcup_{\omega} \mathbb{L}_{\omega}^{1,2}$	$\bigcup_{\omega} \mathbb{L}_{\omega}^0$
(a4)	$\bigcup_{\omega} \mathbb{L}_2 \cup \bigcup_{\omega} \mathbb{L}_2^0 \cup \mathbb{L}_4$	$\bigcup_{\omega} \mathbb{L}_2 \cup \mathbb{L}_2^0 \cup \mathbb{L}_4$	$\mathbb{L}_2 \cup \bigcup_{\omega} \mathbb{L}_2^0 \cup \mathbb{L}_4^0$	$\bigcup_{\omega} \mathbb{L}_2^0 \cup \mathbb{L}_2^1 \cup \mathbb{L}_4^{0,2}$	$\bigcup_{\omega} \mathbb{L}_2 \cup \mathbb{L}_4$
(a5)	$\bigcup_{\omega} \mathbb{L}_2 \cup \bigcup_{\omega} \mathbb{L}_2^0 \cup \mathbb{L}_3$	$\bigcup_{\omega} \mathbb{L}_2 \cup \mathbb{L}_2^0 \cup \mathbb{L}_3$	$\mathbb{L}_2 \cup \bigcup_{\omega} \mathbb{L}_2^0 \cup \mathbb{L}_3$	$\bigcup_{\omega} \mathbb{L}_2^0 \cup \mathbb{L}_2^1 \cup \mathbb{L}_3^0$	$\bigcup_{\omega} \mathbb{L}_2 \cup \mathbb{L}_3$

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Let $\mathbb{X}_i = \langle X_i, \rho_i \rangle$, $i \in I$, be disjoint, connected and reversible L_b -structures. Then the union $\bigcup_{i \in I} \mathbb{X}_i$ is reversible if there is no infinite sequence $\langle i_k : i \in \omega \rangle$ of different elements from I such that $\text{Mono}(\mathbb{X}_{i_{k+1}}, \mathbb{X}_{i_k}) \neq \emptyset$ for each $k \in \omega$, and $\text{Mono}(\mathbb{X}_{i_1}, X_{i_0}) \neq \text{Iso}(\mathbb{X}_{i_1}, \mathbb{X}_{i_0})$.

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If \mathcal{A}_{κ_*} is an infinite family of disjoint, connected, reversible and isomorphic binary structures of size κ_* , such that $\text{Mono}(\mathbb{X}) = \text{Aut}(\mathbb{X})$ for each $\mathbb{X} \in \mathcal{A}_{\kappa_*}$, then $\text{Mono}(\mathbb{X}, \mathbb{Y}) = \text{Iso}(\mathbb{X}, \mathbb{Y})$ for each $\mathbb{X}, \mathbb{Y} \in \mathcal{A}_{\kappa_*}$, $\mathbb{X} \neq \mathbb{Y}$.

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For $k \in \omega$ and $*$ $\in \{\prime, \#, \pm, =, \rightleftharpoons, \Delta\}$, let $\mathbb{Y}_i^{\theta^*=k}$, $i \in M_k^*$, be all finite and finitely many infinite (if there are any) nonisomorphic connected reversible structures for which $\theta^* = k$, and let $\mathbb{Z}_i^{\theta^*=k}$, $i \in N_k^*$, be all nonisomorphic...

Further examples of reversible disconnected structures

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$\theta^*(\mathbb{X})$	$ X $	$ \rho $	$\text{Deg}^\pm(\mathbb{X})$	$ \rho \cap \Delta_X $	$\frac{1}{2} \cdot (\rho \setminus \Delta_X) \cap (\rho \setminus \Delta_X)^{-1} $	\dots

For $k \in \omega$ and $* \in \{\prime, \#, \pm, =, \rightleftharpoons, \Delta\}$, let $\mathbb{Y}_i^{\theta^*=k}$, $i \in M_k^*$, be all finite and finitely many infinite (if there are any) nonisomorphic connected reversible structures for which $\theta^* = k$, and let $\mathbb{Z}_i^{\theta^*=k}$, $i \in N_k^*$, be all nonisomorphic...

- $* \in \{\prime\}$... connected structures for which $\theta' = k$.

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For $k \in \omega$ and $* \in \{!, \#, \pm, =, \rightleftarrows, \Delta\}$, let $\mathbb{Y}_i^{\theta^*=k}$, $i \in M_k^*$, be all finite and finitely many infinite (if there are any) nonisomorphic connected reversible structures for which $\theta^* = k$, and let $\mathbb{Z}_i^{\theta^*=k}$, $i \in N_k^*$, be all nonisomorphic...

- $* \in \{!\}$... connected structures for which $\theta^! = k$.
- $* \in \{\#\}$... digraph trees for which $\theta^\# = k$.

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- $* \in \{\pm\}$... connected (deg^\pm -)regular digraphs for which $\theta^\pm = k$.

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- $* \in \{=\}$... reflexive digraph trees for which $\theta^= = k$.

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- $*$ $\in \{\rightleftarrows\}$... connected graphs for which $\theta^{\rightleftarrows} = k$.

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- $* \in \{=\}$... reflexive digraph trees for which $\theta^= = k$.
- $* \in \{\rightleftarrows\}$... connected graphs for which $\theta^{\rightleftarrows} = k$.
- $* \in \{\Delta\}$... poset trees with no leaves on first level, for which $\theta^\Delta = k$.

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- $* \in \{\rightleftarrows\}$... connected graphs for which $\theta^{\rightleftarrows} = k$.
- $* \in \{\Delta\}$... poset trees with no leaves on first level, for which $\theta^\Delta = k$.

Then all the components $\mathbb{Z}_i^{\theta^*=k}$ are connected, finite and reversible, except the components $\mathbb{Z}_i^{\theta^\pm=k}$ which are at most countable and reversible.

Further examples of reversible disconnected structures

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$*$	k	$ Y_i^{\theta^*=k} $	$ M_k^* $	$ Z_i^{\theta^*=k} $	$ N_k^* $
$/$	$k \in \omega$	$< \omega$	$< \omega$	$< \omega$	$< \omega$
$\#$	$k \in \mathbb{N}$	$< \omega$	$< \omega$	$< \omega$	$< \omega$
\pm	$k \leq 1$	≤ 2	≤ 2	≤ 2	≤ 2
	$k \geq 2$	$\leq \omega$	ω	$\leq \omega$	\mathfrak{c}
$=$	$k \in \omega$	$< \infty$	ω	$< \omega$	$< \omega$
\Leftrightarrow	$k \in \omega$	$< \infty$	ω	$< \omega$	$< \omega$
Δ	$k \in \omega$	$< \infty$	ω	$< \omega$	$< \omega$

Table 2: The size of the sets $Y_i^{\theta^*=k}$, M_k^* , $Z_i^{\theta^*=k}$ and N_k^*

Further examples of reversible disconnected structures

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For $n \geq 2$, and $* \in \{!, \#, \pm, =, \rightleftharpoons, \Delta\}$, let us define the following structure:

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For $n \geq 2$, and $* \in \{!, \#, \pm, =, \rightleftharpoons, \Delta\}$, let us define the following structure:

$$\mathbb{X}_n^* := \bigcup_{k=0}^{n-1} \left(\bigcup_{i \in M_k^*} \mathbb{Y}_i^{\theta^* = k} \right) \cup \bigcup_{i \in N_n^*} \left(\bigcup_{\omega} \mathbb{Z}_i^{\theta^* = n} \right).$$

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All the structures \mathbb{X}_n^* are reversible by the last corollary.

Further examples of reversible disconnected structures

For $n \geq 2$, and $* \in \{!, \#, \pm, =, \leftrightarrow, \Delta\}$, let us define the following structure:

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	\mathbb{X}'_n	$\mathbb{X}_n^\#$	\mathbb{X}_n^\pm	$\mathbb{X}_n^=$	$\mathbb{X}_n^\leftrightarrow$	\mathbb{X}_n^Δ
$n = 2$	(a4) (b2)	(a4) (b2)	(a1) (b2)	(a1) (b2)	(a1) (b2)	(a1) (b2)
$n \geq 3$	(a1) (b2)	(a4) (b2)	(a1) (b2)	(a1) (b2)	(a1) (b2)	(a1) (b2)

Table 3: The place of the structures \mathbb{X}_n^* in Table 1

Further examples of reversible disconnected structures

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	\mathbb{X}'_n	$\mathbb{X}_n^\#$	\mathbb{X}_n^\pm	$\mathbb{X}_n^=$	$\mathbb{X}_n^\leftrightarrow$	\mathbb{X}_n^Δ
$n = 2$	(a4) (b2)	(a4) (b2)	(a1) (b2)	(a1) (b2)	(a1) (b2)	(a1) (b2)
$n \geq 3$	(a1) (b2)	(a4) (b2)	(a1) (b2)	(a1) (b2)	(a1) (b2)	(a1) (b2)

Table 3: The place of the structures \mathbb{X}_n^* in Table 1

Most of the structures \mathbb{X}_n^* are placed in (a1) row of Table 1. The condition (a1) is not operative, therefore the presented sufficient conditions do not substitute each other.

Further examples of reversible disconnected structures

For $n \geq 2$, and $* \in \{!, \#, \pm, =, \leftrightarrow, \Delta\}$, let us define the following structure:

$$\mathbb{X}_n^* := \bigcup_{k=0}^{n-1} \left(\bigcup_{i \in M_k^*} \mathbb{Y}_i^{\theta^*=k} \right) \cup \bigcup_{i \in N_n^*} \left(\bigcup_{\omega} \mathbb{Z}_i^{\theta^*=n} \right).$$

All the structures \mathbb{X}_n^* are reversible by the last corollary.

	\mathbb{X}'_n	$\mathbb{X}_n^\#$	\mathbb{X}_n^\pm	$\mathbb{X}_n^=$	$\mathbb{X}_n^\leftrightarrow$	\mathbb{X}_n^Δ
$n = 2$	(a4) (b2)	(a4) (b2)	(a1) (b2)	(a1) (b2)	(a1) (b2)	(a1) (b2)
$n \geq 3$	(a1) (b2)	(a4) (b2)	(a1) (b2)	(a1) (b2)	(a1) (b2)	(a1) (b2)

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Thank you for your attention.