

Complete pairs of coanalytic sets

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Summary

- 1 Borel, analytic and coanalytic sets
- 2 Definition of a complete pair
- 3 Basic examples of complete pairs
- 4 A complete pair in the space of continuous functions
- 5 A complete pair in the theory of automata



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Definition

X is a **Polish space** if X is separable and completely metrizable.

Cantor set \mathcal{C} , the reals \mathbb{R} , the naturals \mathbb{N} , the Banach space $C([0, 1])$ with $\|\cdot\|_\infty$ are all examples of Polish spaces.



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Definition

The Borel sets $\mathcal{B}(X)$ in a given topological space is the smallest σ -field containing all open sets of X .



Definition

A set $A \subset X$ in a Polish space X is **analytic** if there exists a Polish space Y and a Borel set $B \subset X \times Y$ such that

$$A = \{x \in X : \exists y \in Y \langle x, y \rangle \in B\}.$$

Definition

A set $A \subset X$ in a Polish space X is **coanalytic** if $X \setminus A$ is an analytic set.



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A set A in a topological space X is **Wadge reducible** to a set B in a topological space Y if there exists a continuous mapping $\phi : X \rightarrow Y$ such that $A = \phi^{-1}[B]$.

Definition

A disjoint pair A, B in a topological space X is **Wadge reducible** to a disjoint pair C, D in a topological space Y , if there exists a continuous mapping $\phi : X \rightarrow Y$ such that $A \leq_{\phi} C$ and $B \leq_{\phi} D$, that is $A = \phi^{-1}[C]$ and $B = \phi^{-1}[D]$.



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Definition

A disjoint pair of coanalytic sets C, D in a Polish space X is **complete**, if for every disjoint pair of coanalytic sets A, B in the Cantor set the pair A, B is Wadge reducible to the pair C, D .

The pair C, D represents all essential properties of pairs of coanalytic sets. For example, in the class of coanalytic sets there exists a pair A, B not separable by a Borel set. The same holds for all complete pairs.



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In order to prove that a given disjoint pair C, D of coanalytic sets is complete, it is enough to find a complete pair A, B and a reduction ϕ such that $A \leq_{\phi} C$ and $B \leq_{\phi} D$.



Definition

$T \subset \omega^{<\omega}$ is a **tree**, if T is closed with respect to initial segments, that is for every $s \in T$ and an initial segment $r \preceq s$ we have $r \in T$. A sequence $x \in \omega^\omega$ is a **branch** of T , if for every $n \in \omega$ we have $x|n \in T$.



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Definition

Let $\text{Tr} \subset 2^{\omega^{<\omega}}$ be the set of all trees. We define WF as the set of **all well-founded trees** and UB as the set of **all trees with exactly one branch**.

J. Saint Raymond proved in 2007 that the pair WF, UB is a complete pair of coanalytic sets.



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*Every tree $T \in \text{WF}$ admits a natural **rank** $\text{rk}(T)$, which is an ordinal below ω_1 . Firstly we define inductively rank of T for every vertex of T and then define rank of T as the rank of $\emptyset \in T$. If T is not in WF , we define $\text{rk}(T)$ as ω_1 .*



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Definition

Let

$$V_0 = \{\langle S, T \rangle : S \in \text{WF}, \text{rk}(S) < \text{rk}(T)\}$$

and

$$V_1 = \{\langle S, T \rangle : T \in \text{WF}, \text{rk}(T) \leq \text{rk}(S)\}.$$

The sets V_0 i V_1 are disjoint and coanalytic and forms a complete pair.



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Definition

We define Diff as a subset of $C([0, 1])$ consisting of all **differentiable functions** on the unit interval $[0, 1]$.

In 1936 S. Mazurkiewicz proved that the set Diff is an coanalytic non-Borel subset $C([0, 1])$.

Definition

Let Diff_1 be the set of all functions in $C([0, 1])$ which are **not differentiable in exactly one point** of $[0, 1]$.

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Let \mathcal{S} be the set of all **full binary trees** with vertices **labeled** by elements of the set $\{\exists, \forall\} \times \{0, 1\}$. Let $t \in \mathcal{S}$.

From a vertex of t one may go either right or left and the players \exists and \forall play a **game**, such that each of the players decides about a move from ‘his’ vertices, that is from vertices labeled by \exists and \forall respectively.



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The player \exists wins, if all vertices occurring in a given play, with except of finitely many, have label 0. **The player \forall wins**, if all vertices occurring in a given play, with except of finitely many, has label 1.



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Definition

Let $W_{0,1}$ be the set of all trees in \mathcal{S} , such that the **player \exists has a winning strategy** and $W'_{0,1}$ be the set of all trees in \mathcal{S} , such that the **player \forall has a winning strategy**.

The pair $W_{0,1}, W'_{0,1}$ is a complete pair of coanalytic sets. Sz. Hummel proved in his Master Dissertation that the sets $W_{0,1}, W'_{0,1}$ are coanalytic and that the sets $W_{0,1}, W'_{0,1}$ are not separable by a Borel set. This results were incorporated into a joint paper by D. Niwiński, Sz. Hummel and H. Michalewski accepted for STACS 2009.



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The set $W_{0,1}$ is an example of set accepted by a **non-deterministic tree automaton**. The automaton has the **states** $0, 1$ and T , works over the **alphabet** $\{\exists, \forall\} \times \{0, 1\}$ and has the following **transitions**:

$$i \xrightarrow{\langle \forall, j \rangle} j, j,$$

$$i \xrightarrow{\langle \exists, j \rangle} j, T, \quad i \xrightarrow{\langle \exists, j \rangle} T, j$$

and

$$T \xrightarrow{a} T, T,$$

where $i, j \in \{0, 1\}$ and $a \in \{\exists, \forall\} \times \{0, 1\}$.



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The **rank** of the states 0 and T is 0 and the rank of the state 1 is 1. A tree $t \in \mathcal{S}$ is **recognized** by the automaton if there exists a run of the automaton such that on every branch x of t the $\limsup \rho(x(n))$ is even (in our case the only possible even rank is 0). The set $W'_{0,1}$ is accepted by a very similar automaton, such that the roles of \exists and \forall are swapped and at the same time the roles of 0 and 1 are swapped.



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An analogous definition gives sets $W_{i,k}, W'_{i,k}$ for larger sets of indices $\{i, \dots, n\}$. One can prove, that the **complement of the set $W_{0,1}$ is not recognized** by an automaton with index $\{0, 1\}$ but is accepted by an automaton with index $\{1, 2\}$.



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It was proved by M. O. Rabin, that every two disjoint sets of trees A , B accepted by automata with indices $\{1, 2\}$ is possible to separate by set C , which is simultaneously accepted by automata with indices $\{1, 2\}$ and $\{0, 1\}$.

For larger set of indices the question of separation remains open.



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Bibliography I



Sz. Hummel

Własności oddzielania zbiorów drzew definiowalnych przez automaty

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J. Saint Raymond

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