Big Ramsey degrees of homogeneous structures part 1: the prehistory

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I would like to cover some old&new results in the area of big Ramsey degrees:

- The prehistory: From Ramsey theorem to big Ramsey degrees of the order of rationals and of the countable random graph.
 (Milliken tree theorem, envelopes and embedding types, optimality).
- Big Ramsey degrees using parameter spaces: Recent results on triangle-free graphs, posets and metric spaces using the Carlson-Simpson theorem.
- Higher arities and bigger forbidden substructures: Hypergraphs, free amalgamation classes.

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 $N \longrightarrow (n)_{k,t}^{p}$: For every partition of $\binom{\omega}{p}$ into *k* classes (colours) there exists $X \in \binom{\omega}{\omega}$ such that $\binom{X}{p}$ belongs to at most *t* parts.

 $(t = 1 \text{ means that } \begin{pmatrix} \chi \\ \rho \end{pmatrix}$ is monochromatic.)

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In 1970's a concept of structural Ramsey theory was introduced. A Ramsey theorem can be seen as a theorem about the class of linear orders.

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Let \mathcal{O} be the class of all finite linear orders.

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 $\begin{pmatrix} B \\ A \end{pmatrix}$ is the set of all embeddings of structure **A** to structure **B**.

 $C \longrightarrow (B)_{k,t}^{A}$: For every *k*-colouring of $\binom{C}{A}$ there exists $f \in \binom{C}{B}$ such that $\binom{f(B)}{A}$ has at most *t* colours.

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A natural question: Is the same true for (\mathbb{Q}, \leq) (the order of rationals)?

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Sierpiński: not true for |O| = 2.

























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$$\forall_{(\mathcal{O},\leq_{\mathcal{O}})\in\mathcal{O}}\exists\tau=\tau(|\mathcal{O}|)\in\omega}\forall_{k\geq1}:(\mathbb{Q},\leq)\longrightarrow(\mathbb{Q},\leq)_{k,T}^{(\mathcal{O},\leq_{\mathcal{O}})}.$$

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Trees (terminology)

- A tree is a (possibly empty) partially ordered set (*T*, <_T) such that, for every *t* ∈ *T*, the set { *s* ∈ *T* : *s* <_T *t* } is finite and linearly ordered by <_T. All trees considered are finite or countable.
- All nonempty trees we consider are rooted, that is, they have a unique minimal element called the root of the tree.
- An element t ∈ T of a tree T is called a node of T and its level, denoted by |t|_T, is the size of the set { s ∈ T : s <_T t }.
- We use *T*(*n*) to denote the set of all nodes of *T* at level *n*,
- For $s, t \in T$, the meet $s \wedge_T t$ of s and t is the largest $s' \in T$ such that $s' \leq_T s$ and $s' \leq_T t$.
- The height of *T*, denoted by h(T), is the minimal natural number *h* such that $T(h) = \emptyset$. If there is no such number *h*, then we say that the height of *T* is ω .



Subtrees and strong subtrees

- A subtree of a tree *T* is a subset *T'* of *T* viewed as a tree equipped with the induced partial ordering such that s ∧_{T'} t = s ∧_T t for each s, t ∈ *T'*.
- Given a tree T and nodes $s, t \in T$ we say that s is a successor of t in T if $t \leq_T s$.
- The node s is an immediate successor of t in T if t <_T s and there is no s' ∈ T such that t <_T s' <_T s.
- We denote the set of all successors of t in T by Succ_T(t) and the set of immediate successors of t in T by ImmSucc_T(t).

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Definition (Strong subtree)

A subtree S of a tree T is a strong subtree of T if either S is empty, or S is nonempty and satisfies the following three conditions.

- 1 The tree *S* is rooted and balanced.
- **2** Every level of *S* is a subset of some level of *T*, that is, for every n < h(S) there exists $m \in \omega$ such that $S(n) \subseteq T(m)$.
- G For every non-maximal node s ∈ S and every t ∈ ImmSucc_T(s) the set ImmSucc_S(s) ∩ Succ_T(t) is a singleton.

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3 Every level of *S* is a subset of some level of *T*.

4 *S* either has no leaves or all are at the same level.

Let *T* be a tree and $k \in \omega + 1$. We use $Str_k(T)$ to denote the set of all strong subtrees of *T* of height *k*.

Theorem (Milliken 1979)

For every rooted, balanced and finitely branching tree T of infinite height, every $k \in \omega$ and every finite colouring of $\operatorname{Str}_k(T)$ there is $S \in \operatorname{Str}_\omega(T)$ such that the set $\operatorname{Str}_k(S)$ is monochromatic.

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Notice that for regularly branching tree the strong subtree is isomorphic to the original tree.

We aim to prove:

Theorem (Laver, late 1969)

$$\forall_{(O,\leq_O)\in\mathcal{O}}\exists_{\mathcal{T}=\mathcal{T}(|O|)\in\omega}\forall_{k\geq 1}: (\mathbb{Q},\leq)\longrightarrow (\mathbb{Q},\leq)_{k,\mathcal{T}}^{(O,\leq_O)}.$$

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 Nodes of the binary tree 2^{<ω} ordered "from left to right" yields (Q, ≤).



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- Using the Halpern–Läuchli's theorem we can find a strong subtree which is monochromatic.
- ④ Strong subtree is isomorphic to the original binary tree and also isomorphic to (Q, ≤) when ordered lexicographically
 - \implies we found the monochromatic copy!



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If |O| = n > 1 we transfer colourings of *n*-tuples of nodes to colouring of strong subtrees.





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Multiple choices of X may lead to a same envelope. We speak of different embedding types within a given envelope.

Now we can finish proof of:

Theorem (Laver, late 1969)

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Proof.

• Fix $(O, \leq_O) \in \mathcal{O}$ and put n = |O|.

2 T(n) is the number of embedding types of *n*-tuples in the binary tree.

- Recall that height of each envelope is at most 2n 1.
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- For each embedding type construct colouring of envelopes and pass to a monochromatic subtree by the application of Milliken tree theorem.
- **5** The resulting copy will have at most T(n) different colours

Big Ramsey degrees of (\mathbb{Q}, \leq) are finite!



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Fun fact: Number of Devlin types of size n is

$$t_n = \sum_{\ell=1}^{n-1} {\binom{2n-2}{2\ell-1}} t_\ell \cdot t_{n-1}$$
 with $n_1 = 1$

This is well known sequence (of the odd tangent numbers).

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Where minimal T satisfying the statement above is the number of Devlin types of size |O|

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Proof.



Theorem (Devlin, 1979)

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If $B \subset A$ is a subset of Devlin type A, then the embedding type of B inside every minimal envelope is a Devlin type.

Proof.

Recall that in Devlin type on every is either leaf or branching. Minimal envelope will include all those levels where branching or leaf of *B* happens and skip all others.

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Proof.

Recall that in Devlin type on every is either leaf or branching. Minimal envelope will include all those levels where branching or leaf of *B* happens and skip all others.

We thus obtain an upper bound: T(|O|) is at most the number of Devlin types.

Lemma

Let A be a Devlin type representing (\mathbb{Q}, \leq) and $B \subseteq A$ a copy of (\mathbb{Q}, \leq) then there exists a $C \subseteq B$ whose embedding type (inside a minimal envelope) is A.

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Proceed by induction on the individual branching/leaf events of A.



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Proceed by induction on the individual branching/leaf events of *A*. First produce meet of *B*. *A* starts with a branching. Place the root of *A* to this meet.



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Proof by a slide-show.

Proceed by induction on the individual branching/leaf events of *A*. The meet splits leafs of *B* into two infinite intervals. Each with meet above its son.



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Proof by a slide-show.

Proceed by induction on the individual branching/leaf events of *A*. Use corresponding meet to realize 2nd branching of *A*.



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Proof by a slide-show.

Proceed by induction on the individual branching/leaf events of *A*. *B* further subdivides into intervals.



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Proof by a slide-show.

Proceed by induction on the individual branching/leaf events of *A*. Continue analogously with next branching.



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Proof by a slide-show.

Proceed by induction on the individual branching/leaf events of *A*. Place the first leaf of *A* into corresponding interval.



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Proceed by induction on the individual branching/leaf events of *A*. ... blablabla... proof done.



Victory!



We characterised the big Ramsey degrees of rationals and gave a closed-form formula. This shows that the upper bound proof is best possible. It also has applications to topological dynamics

Zucker: Big Ramsey degrees and topological dynamics

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- We denote by \mathcal{G} the class of all finite graphs.

Theorem

$$\forall_{\mathbf{A}\in\mathcal{G}}\exists_{\mathcal{T}=\mathcal{T}'(\mathbf{A})\in\omega}\forall_{k\geq 1}:\mathbf{R}\longrightarrow(\mathbf{R})_{k,\mathcal{T}}^{\mathbf{A}}.$$

This theorem was published by Sauer in 2006 and also appears in Todorčević' Introduction to Ramsey spaces. Values of $T'(\mathbf{G})$ were characterised by Laflamme–Sauer–Vuksanović in 2010.

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A finitary version is (probably more) famous!

Theorem (Nešetřil-Rödl 1977, Abramson-Harington 1978)

$$\forall_{\mathbf{A}\in\mathcal{G}}\exists_{t=t(\mathbf{A})\in\omega}\forall_{\mathbf{B}\in\mathcal{G},k\geq 1}\exists_{\mathbf{C}}\in\mathcal{G}:\mathbf{C}\longrightarrow(\mathbf{B})_{k,t}^{\mathbf{A}}.$$

Understanding the unavoidable colourings

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For (\mathbb{Q}, \leq) we have the Sierpiński colourings. Can we do something similar for the Rado graph?



















Definition (Graph G)

- 1 Vertices: $2^{<\omega}$
- 2 Vertices $a, b \in 2^{<\omega}$ satisfying |a| < |b| forms and edge if and only if b(|a|) = 1.
- **3** There are no other edges.



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The upper bound

Lemma

G is universal: the Rado graph R embeds to G.

Proof.

Assume that the vertex set of **R** is ω . The vertex $i \in \omega$ then corresponds to a sequence *a* of length *i* with a(j) = 1 if and only if $i \sim j$.

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The definition of **G** is stable for passing into a strong subtrees: if *S* is a strong subtree of $2^{<\omega}$ then it is also a copy of **G** in **G**

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Lower bounds needs a bit more care.

Thank you for the attention

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(See also S. Todorčević, Introduction to Ramsey spaces.)